## Graphs

## Data Structures: Review

- Sequential Data Structures

| Index | 0 | 1 | 2 |
| :---: | :---: | :---: | :---: |
| Value | 5 | 20 | 3 |
| Node | $\rightarrow$ | Node | Node |
| 5 |  | 20 | 3 |

- Key-Value Store
- Stores pairs of elements with no particular order
- Each key is associated with one value

| Key | "cse" | "mga" | "geo" |
| :---: | :---: | :---: | :---: |
| Value | 20 | 3 | 5 |

- Ex. Map, Dictionary, Object
- Tree
- Non-linear structure
- Each element can be associated with multiple other elements



## Data Structures

- How do we store data with multiple interconnected associations?
- A [station, intersection, city] can have multiple connections



## Data Structures

- Let's use trees
- Start with UCLA as the root
- Recursively add all children

- Oops
- We have duplicates in our data structure


## Data Structures

- Let's try again
- When we try to add a duplicate, add a reference to the existing node



## Graphs

- This is a graph
- Similar to a tree, except cycles are allowed
- Cycle: Can "travel" from a node back to itself without repeating a node



## Graphs

- Because of the cycles, our tree traversals will get stuck in infinite recursion
- No leaves (node with no children) to terminate the recursion



## Graphs

- We'll need a new way of representing this data structure and new algorithms to work with the data
- Store the nodes and edges



## Graphs - Nodes and Edges

- Node: Each data element is stored in a node, similar to linked lists and trees
- Edge: A connection between two nodes



## Graphs - Adjacency List

- A map of nodes to all nodes connected to it through an edge
- This is how we'll represent graphs


| UCLA | STANFORD, SRI, UCSB, RAND |
| :---: | :---: |
| STANFORD | UCLA, SRI |
| SRI | UCLA, UCSB, UTAH, STANFORD |
| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC, BBN |
| UTAH | SRI, SDC, MIT |
| SDC | UTAH, RAND |
| MIT | UTAH, BBN, LINCOLN |
| BBN | RAND, MIT, HAVARD |
| LINCOLN | MIT, CASE |
| CARNEGIE | CASE, HARVARD |
| HARVARD | BBN, CARNEGIE |
| CASE | LINCOLN, CARNEGIE |

## Graphs - Adjacency List

- We use generics in our code so we can create a graph of any type

```
public class Graph<N> {
private HashMap<N, ArrayList<N>> adjacencyList;
public Graph() {
    this.adjacencyList = new HashMap<>();
}
public void addEdge(N from, N to) {
    this.addNode(from);
    this.addNode(to);
    this.adjacencyList.get(from).add(to);
}
public void addBidirectionalEdge(N node1, N node2) {
    this.addNode(node1);
    this.addNode(node2);
    this.adjacencyList.get(node1).add(node2);
    this.adjacencyList.get(node2).add(node1);
}
private void addNode(N a) {
    if (!this.adjacencyList.containsKey(a)) {
        this.adjacencyList.put(a, new ArrayList<>());
        }
```

$\}$

```
public static void example(){
Graph<String> graph = new Graph<>();
graph.addBidirectionalEdge("UCLA","STANFORD");
graph.addBidirectionalEdge("UCLA","SRI");
graph.addBidirectionalEdge("UCLA","UCSB");
graph.addBidirectionalEdge("UCLA","RAND");
graph.addBidirectionalEdge("STANFORD","SRI");
graph.addBidirectionalEdge("SRI","UCSB");
graph.addBidirectionalEdge("UCSB","RAND");
graph.addBidirectionalEdge("SRI","UTAH");
graph.addBidirectionalEdge("RAND","SDC");
graph.addBidirectionalEdge("UTAH","SDC");
graph.addBidirectionalEdge("UTAH", "MIT");
graph.addBidirectionalEdge("RAND","BBN");
graph.addBidirectionalEdge("MIT","BBN");
graph.addBidirectionalEdge("MIT","LINCOLN");
graph.addBidirectionalEdge("LINCOLN","CASE");
graph.addBidirectionalEdge("CASE","CARNEGIE");
graph.addBidirectionalEdge("CARNEGIE", "HARVARD");
graph.addBidirectionalEdge("HARVARD","BBN");
```


## Paths

- A path is a sequence of nodes where each pair of adjacent nodes are connected by an edge
- ["UCLA", "SRI", "UTAH", "MIT", "BBN", "RAND"] is a path in this graph
- ["SRI", "UTAH", "BBN"] is not a path since UTAH and BBN are not connected by an edge


| UCLA | STANFORD, SRI, UCSB, RAND |
| :---: | :---: |
| STANFORD | UCLA, SRI |
| SRI | UCLA, UCSB, UTAH, STANFORD |
| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC, BBN |
| UTAH | SRI, SDC, MIT |
| SDC | UTAH, RAND |
| MIT | UTAH, BBN, LINCOLN |
| BBN | RAND, MIT, HAVARD |
| LINCOLN | MIT, CASE |
| CARNEGIE | CASE, HARVARD |
| HARVARD | BBN, CARNEGIE |
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## Breadth-First Search (BFS)

## Connected Component

- This graph is connected
- There exists a path between any 2 nodes in the graph


| UCLA | STANFORD, SRI, UCSB, RAND |
| :---: | :---: |
| STANFORD | UCLA, SRI |
| SRI | UCLA, UCSB, UTAH, STANFORD |
| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC, BBN |
| UTAH | SRI, SDC, MIT |
| SDC | UTAH, RAND |
| MIT | UTAH, BBN, LINCOLN |
| BBN | RAND, MIT, HAVARD |
| LINCOLN | MIT, CASE |
| CARNEGIE | CASE, HARVARD |
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## Connected Component

- What if a few connections are broken?
- How can we tell if two nodes are connected?


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| STANFORD | UCLA, SRI |
| SRI | UCLA, UCSB, UTAH, STANFORD |
| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC |
| UTAH | SRI, SDC |
| SDC | UTAH, RAND |
| MIT | BBN, LINCOLN |
| BBN | MIT, HAVARD |
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## Connected Component

- We could verify manually for this graph
- But the Internet has gotten a little bigger over time
- Need to code an algorithm to solve this for us



## BFS

- The Algorithm: Breadth-First Search (BFS)
- Choose a starting node
- Continuously explore connected nodes


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| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC |
| UTAH | SRI, SDC |
| SDC | UTAH, RAND |
| MIT | BBN, LINCOLN |
| BBN | MIT, HAVARD |
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## BFS

- Choose a starting node


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## BFS

- Explore all nodes connected to the striating node


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| UCSB | UCLA, SRI, RAND |
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## BFS

- Repeatedly explore nodes that were visited in the last round


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| UCSB | UCLA, SRI, RAND |
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## BFS

- Repeat until no new nodes are added
- Never visit a node twice


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| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC |
| UTAH | SRI, SDC |
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## BFS

- Use a queue to track the order of nodes to visit
- Start with starting node in the queue
- When visiting a node, add all unexplored neighbors to the queue
- Visit neighbors of the node at the front of the

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| UTAH | SRI, SDC |
| SDC | UTAH, RAND |
| MIT | BBN, LINCOLN |
| BBN | MIT, HAVARD |
| LINCOLN | MIT, CASE |
| CARNEGIE | CASE, HARVARD |
| HARVARD | BBN, CARNEGIE |
| CASE | LINCOLN, CARNEGIE | queue until the queue is empty

## BFS

- More BFS details to come

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| :---: | :---: |
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| UCSB | UCLA, SRI, RAND |
| RAND | UCLA, UCSB, SDC |
| UTAH | SRI, SDC |
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| BBN | MIT, HAVARD |
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## Connectivity

- If you start at nodeA and explore nodeB during the algorithm
- nodeA and nodeB are connected

